

## $n$ -Tuple Colorings and Associated Graphs

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The  $n$ -tuple graph coloring, which assigns to each vertex  $n$  colors, is defined together with its respective chromatic number  $\chi_n$ . It is proved that these numbers satisfy the inequality  $\chi_n \geq 2 + \chi_{n-1}$ , and that equality holds only for bipartite graphs. Graphs  $G_n^m$  are defined which play the same role for the  $n$ -tuple coloring that  $K_m$  plays for the conventional coloring. The chromatic numbers of various classes of graphs are also calculated.

### I. DEFINITIONS AND INTRODUCTION

A graph  $G$  is understood to be finite and to contain no loops or multiple edges. Its vertex set is denoted by  $V(G)$ .  $C_n$  and  $K_n$  represent, respectively, the  $n$ -cycle and the complete graph on  $n$  vertices. The cardinality of a set  $S$  is denoted by  $|S|$ .

$I^+$  is the set of all positive integers and  $I^m$  is the set of all positive integers not greater than  $m$ .  $I_n^m$  is the family of all subsets of  $I^m$  of cardinality  $n$ . We set the convention that when any member of  $I_n^m$  is written out explicitly, its elements should be listed in monotone increasing order. The graph  $G_n^m$  is now defined as follows,  $V(G_n^m) = I_n^m$  and two vertices are adjacent iff they are disjoint. For example, Fig. 1 is a drawing of  $G_2^5$ .

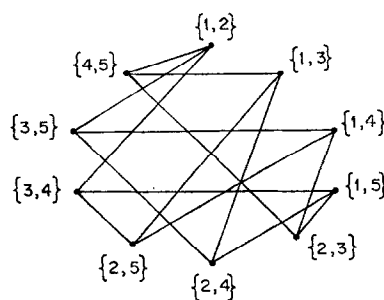


FIGURE 1

These graphs were originally defined and to some extent discussed in [2]. If  $m < 2n$  then  $G_n^m$  has no edges and is of no interest to us. Consequently, whenever a graph  $G_n^m$  is considered in the sequel, it is implicitly assumed that  $m \geq 2n$ . If  $n = 1$ , then any two vertices are adjacent and  $G_1^m = K_m$ .

An  $n$ -tuple coloring of a graph  $G$  is an assignment of  $n$  distinct colors to each vertex of  $G$  in such a manner that no two adjacent vertices share a color. This is clearly a generalization of the conventional coloring, which assigns only one color to each vertex. The  $n$ -tuple coloring can be regarded from two different points of view, and both will be used frequently in this paper. The first of these makes use of the notion of homomorphism. We recall that a map  $\eta: V(G) \rightarrow V(H)$  is said to be a *homomorphism of  $G$  into  $H$*  if  $\eta$  maps pairs of adjacent vertices of  $G$  onto pairs of adjacent vertices of  $H$ . Now suppose that an  $n$ -tuple coloring of  $G$  utilizes a total of  $m$  colors. We may then denote the colors used in this coloring by  $1, 2, \dots, m$ . The color assignment then becomes a mapping  $\eta: V(G) \rightarrow I_n^m$ . But  $I_n^m$  is the vertex set of  $G_n^m$  and the condition that no two adjacent vertices share a color makes  $\eta$  into a homomorphism of  $G$  into  $G_n^m$ . To see this the reader should recall that two vertices of  $G_n^m$  are adjacent iff they are disjoint. Thus, we have obtained our first reformulation of the concept of an  $n$ -tuple coloring:

AF1. *An  $n$ -tuple coloring of  $G$  with  $m$  colors is a homomorphism  $\eta: G \rightarrow G_n^m$ .*

This alternate formulation will be referred to as AF1.

The second point of view utilizes independent sets. A subset of  $V(G)$  is said to be *independent* if it contains no adjacent vertices. We again identify the colors used with the integers of  $I^m$ . Let  $C_i$  be the set of vertices to which the "color"  $i$  was assigned. The condition that no two adjacent vertices share a color now means that each  $C_i$  is an independent set. Thus,

AF2. *An  $n$ -tuple coloring of  $G$  with  $m$  colors is a collection  $\{C_i\}$  ( $1 \leq i \leq m$ ) of independent subsets of  $V(G)$  such that every vertex of  $G$  is contained in exactly  $n$  of the  $C_i$ .*

As was defined in [2], such a family of sets is said to *cover  $G$  exactly  $n$  times*. This point of view will be referred to as AF2. It should be noted that for  $n = 1$  these points of view are already known and widely used in the literature (see [3, Chap. 12]).

$\chi_n(G)$ , the  $n$ th chromatic number of  $G$ , is the smallest number of colors needed to give  $G$  an  $n$ -tuple coloring. In view of AF1,  $\chi_n(G)$  is the smallest integer  $m$  for which there exists a homomorphism  $G \rightarrow G_n^m$ .

Alternately, AF2 tells us that  $\chi_n(G)$  is the smallest number of independent subsets of  $V(G)$  which cover  $G$  exactly  $n$  times.

Obviously  $\chi_1(G) = \chi(G)$ , the conventional chromatic number. The higher chromatic numbers can also be given an alternate interpretation in terms of  $\chi(G)$ . If  $G[H]$  is the lexicographic product of the graphs  $G$  and  $H$  (see [2]), then an  $n$ -tuple coloring of  $G$  is equivalent to a conventional coloring of  $G[K_n]$ . Therefore,

$$\chi_n(G) = \chi(G[K_n]).$$

This allows us to reinterpret [2, Theorem 3 and the corollary to Theorem 5] as

**THEOREM 5'.** *If  $G$  is critical and not complete then*

$$\chi_n(G) \leq n\chi(G) - (n/2).$$

**THEOREM 3'.**  $\chi(G[H]) = \chi_n(G)$  where  $n = \chi(H)$ .

As these theorems will not be used in this paper, no new proof is offered.

## 2. GENERAL PROPERTIES OF $\chi_n$

Lemma 3 of [2] can be reinterpreted as Theorem 1 below. As this theorem will be used repeatedly in the sequel, we offer a new proof.

**THEOREM 1.**  $\chi_{qn+r}(G) \leq q\chi_n(G) + \chi_r(G)$ .

*Proof.* Let  $\{C_i\}$  and  $\{C'_j\}$  be an  $m$ -tuple and an  $n$ -tuple coloring of  $G$  (see AF2).  $\{C_i\} \cup \{C'_j\}$  is then an  $(m + n)$ -tuple coloring of  $G$ . Therefore, we have

$$\chi_{m+n}(G) \leq \chi_m(G) + \chi_n(G),$$

and the full statement of the theorem follows easily.

Q.E.D.

Theorems 6 and 4 of [2] are both sharpened by the following very useful theorem.

**THEOREM 2.** *If  $G$  has an edge, then  $\chi_n(G) \geq 2 + \chi_{n-1}(G)$  for all  $n > 1$ .*

The proof of this theorem relies on the following lemma.

**LEMMA.** *If  $m > 2n$ , then there exists a homomorphism  $\eta: G_n^m \rightarrow G_{n-1}^{m-2}$ .*

*Proof.* A vertex  $A = \{a_1, a_2, \dots, a_n\}$  of  $G_n^m$  is said to be  $k$ -regular if for some  $k$ ,  $a_k = k$  and  $a_{k+1} > k + 1$ . If a vertex is not  $k$ -regular for any  $k$ , it is said to be irregular. Since we have set the convention that  $a_i < a_{i+1}$  for  $1 \leq i < n$ , it follows that in a  $k$ -regular vertex  $a_i = i$  for all  $i \leq k$  and  $a_i > i$  for all other  $i$ . On the other hand, if  $A$  is irregular then  $a_i > i$  for all  $i \leq n$ . We now defined the map  $\eta$  as follows. If  $A$  is  $k$ -regular then

$$\eta(A) = \{a_2 - 1, a_3 - 1, \dots, a_k - 1, a_{k+1} - 2, a_{k+2} - 2, \dots, a_n - 2\};$$

if  $A$  is irregular then

$$\eta(A) = \{a_2 - 2, a_3 - 2, \dots, a_n - 2\}.$$

It is easily seen that  $\eta$  is in fact a mapping into  $I_{n-1}^{m-2}$ . Hence, it now remains to show that  $\eta$  is in fact a homomorphism of  $G_n^m$  into  $G_{n-1}^{m-2}$ . We must show that if  $A = \{a_1, a_2, \dots, a_n\}$  and  $B = \{b_1, b_2, \dots, b_n\}$  are adjacent vertices of  $G_n^m$  then  $\eta(A)$  and  $\eta(B)$  are adjacent vertices of  $G_{n-1}^{m-2}$ . In other words,  $A \cap B = \emptyset \Rightarrow \eta(A) \cap \eta(B) = \emptyset$ . So now suppose  $A$  and  $B$  are disjoint. Since any regular vertex must necessarily contain the integer 1, it follows that not both  $A$  and  $B$  are regular. The proof is therefore reduced to two cases.

*Case 1.* Both  $A$  and  $B$  are irregular. Then  $\eta(A) = \{a_1 - 2, a_3 - 2, \dots, a_n - 2\}$  and  $\eta(B) = \{b_2 - 2, b_3 - 2, \dots, b_n - 2\}$ . So, if  $\eta(A) \cap \eta(B) \neq \emptyset$ , then there exists some  $a_i$  and  $b_j$  such that  $a_i - 2 = b_j - 2$ , or  $a_i = b_j$ , contradicting the disjointness of  $A$  and  $B$ .

*Case 2.*  $A$  is  $k$ -regular and  $B$  is irregular. Then,

$$\eta(A) = \{a_2 - 1, a_3 - 1, \dots, a_k - 1, a_{k+1} - 2, a_{k+2} - 2, \dots, a_n - 2\}$$

and

$$\eta(B) = \{b_2 - 2, b_3 - 2, \dots, b_n - 2\}.$$

If  $\eta(A) \cap \eta(B) \neq \emptyset$ , then

$$\text{for some } j > 1, \quad b_j - 2 \in \eta(A). \quad (1)$$

If  $b_j - 2 = a_i - 2$  for some  $i > k$ , then  $b_j = a_i$ , contradicting the disjointness of  $A$  and  $B$ . Hence, we must have that  $b_j - 2 = a_i - 1$  for some  $2 \leq i \leq k$ .

$$\therefore b_j = a_i + 1 = i + 1 \leq k + 1. \quad (2)$$

However, since  $A$  is  $k$ -regular it follows that  $I^k \subseteq A$ . Since  $A \cap B = \emptyset$ , this implies that

$$b_s \geq k + 1 \quad \text{for all } b_s \in B. \quad (3)$$

Combining (3) with (2) we obtain:  $b_j = k + 1$ .  $b_j$  is therefore the smallest element of  $B$ , i.e.,  $j = 1$ . This, however, contradicts (1).

Thus, we have derived a contradiction from the assumption that  $\eta(A) \cap \eta(B) \neq \emptyset$ , and, consequently,  $\eta$  preserves adjacency and is a homomorphism. Q.E.D.

*Proof of Theorem 2.* Let  $uv$  be an edge of  $G$  and let  $m = \chi_n(G)$ . In view of AF1, there exists a homomorphism  $\gamma: G \rightarrow G_n^m$ . Since  $u$  and  $v$  are adjacent,  $\gamma(u) \cap \gamma(v) = \emptyset$ , and therefore,  $|\gamma(u) \cup \gamma(v)| = 2n$ . Consequently,  $m \geq 2n$ , and by the lemma there exists a homomorphism  $\eta: G_n^m \rightarrow G_{n-1}^{m-2}$ . Since the composition of two homomorphisms is also a homomorphism it follows that  $\eta \circ \gamma$  is a homomorphism of  $G$  into  $G_{n-1}^{m-2}$ .

$$\therefore \chi_{n-1}(G) \leq m - 2 = \chi_n(G) - 2.$$

Q.E.D.

**COROLLARY.** *If  $G$  has an edge then*

$$\chi_n(G) \geq \chi_k(G) + 2n - 2k \quad \text{if } n \geq k.$$

The following theorem is a natural generalization of a well-known theorem on the chromatic number.

**THEOREM 3.** *If  $\gamma: G \rightarrow H$  is a homomorphism then*

$$\chi_n(G) \leq \chi_n(H) \quad \text{for all } n \geq 1.$$

*Proof.* Let  $m = \chi_n(H)$ ; then there exists a homomorphism  $\beta: H \rightarrow G_n^m$ .

$$\therefore \beta \circ \gamma \text{ is a homomorphism of } G \text{ into } G_n^m.$$

$$\therefore \chi_n(G) \leq m = \chi_n(H).$$

Q.E.D.

**COROLLARY.** *If the graphs  $G$  and  $H$  have homomorphisms  $\alpha: G \rightarrow H$  and  $\beta: H \rightarrow G$ , then  $\chi_n(G) = \chi_n(H)$  for all  $n \geq 1$ .*

The converse to the above corollary does not hold. We shall later see that  $\chi_n(G_2^5) = 2n + 1 + [(n - 1)/2]$ . Moreover, the example near the end of Section 3 introduces another graph  $G$  with the same chromatic numbers. However, the reader can easily verify that neither  $G$  nor  $G_2^5$

contain triangles and yet any homomorphic image of either must contain a triangle. It follows that neither  $G$  nor  $G_2^5$  contains a homomorphic image of the other.

### 3. COMPLETE $n$ -PARTITE AND BIPARTITE GRAPHS, $C_n$ AND $G_n^{2n+1}$

We are now in position to calculate the  $n$ -tuple chromatic numbers of the graphs listed in this chapter's title. A graph is said to be *complete  $n$ -partite* if it has a family  $\{U_i\}$  ( $1 \leq i \leq n$ ) of pairwise disjoint nonvoid sets of vertices such that  $\bigcup_{i=1}^n U_i = V(G)$  and two vertices are adjacent iff they belong to different  $U_i$ . Clearly these  $U_i$  are the maximal independent subsets of  $G$ .

**THEOREM 4.** *If  $G$  is complete  $n$ -partite then  $\chi_k(G) = kn$ .*

*Proof.* Clearly  $\chi_1(G) = n$ , and applying Theorem 1 we get

$$\chi_k(G) \leq kn. \quad (4)$$

On the other hand, let  $\sigma = \{S_\lambda\}_{1 \leq \lambda \leq n}$  be any  $k$ -tuple coloring of  $G$ . For each  $1 \leq i \leq n$ , let  $u_i \in U_i$  and let  $\sigma_i$  be the subfamily of  $\sigma$  which consists of all  $S_\lambda$  containing  $u_i$ . By AF2,  $|\sigma_i| = k$ . Moreover, if  $S_\lambda \in \sigma_i$ , then  $S_\lambda \subseteq U_i$ , and since the  $U_i$  are pairwise disjoint it follows that also  $\sigma_i \cap \sigma_j = \emptyset$  for  $i \neq j$ .

$$\therefore m = |\sigma| \geq \sum_{i=1}^n |\sigma_i| = kn.$$

$$\therefore \chi_k(G) \geq kn.$$

Combining this with (4) we get  $\chi_k(G) = kn$ .

Q.E.D.

A graph  $G$  is said to be *bipartite* if  $\chi(G) = 2$ . We shall later see that bipartite graphs are the only graphs to satisfy the extreme case in Theorem 2. For the moment, however, we must content ourselves with the calculation of their chromatic numbers.

**THEOREM 5.** *If  $G$  is bipartite then  $\chi_n(G) = 2n$ .*

*Proof.* From Theorem 1 it follows that  $\chi_n(G) \leq n\chi_1(G) = 2n$ . On the other hand, it follows from the corollary to Theorem 2 that

$$\chi_n(G) \geq \chi_1(G) + 2n - 2 = 2 + 2n - 2 = 2n.$$

$$\therefore \chi_n(G) = 2n.$$

Q.E.D.

Both trees and cycles of even length are bipartite, so that Theorem 5 gives us all of their chromatic numbers. The derivation of the chromatic numbers of the odd cycles is not quite so easy and requires some preliminary discussions and lemmas.

It is well known that  $\chi_1(C_{2p+1}) = 3$ . Figures 2 and 3 show that  $\chi_2(C_7) \leq 5$  and  $\chi_3(C_7) \leq 7$ . Theorem 2 enables us to change these into the equations  $\chi_2(C_7) = 5$  and  $\chi_3(C_7) = 7$ . On the other hand, a 4-tuple coloring of  $C_7$  with nine colors can not be found so that  $\chi_4(C_7) \geq 10$ . Applying Theorem 1 we get  $\chi_4(C_7) \leq \chi_3(C_7) + \chi_1(C_7) = 7 + 3 = 10$ . We are thus led to conjecture that  $\chi_k(C_{2p+1}) = 2k + 1$  if  $k \leq p$  and  $\chi_{p+1}(C_{2p+1}) = 2p + 4$ .

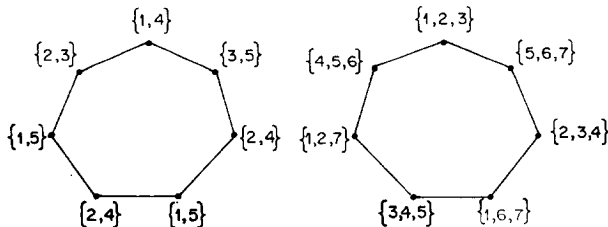
We know that an  $n$ -tuple coloring of  $C_{2p+1}$  with  $m$  colors is equivalent to a homomorphism  $\theta: C_{2p+1} \rightarrow G_n^m$ . As  $\chi_1(C_{2p+1}) = 3$ , we must also have  $\chi_1(\theta(C_{2p+1})) \geq 3$ . By a well-known theorem (see [3, p. 127]),  $\theta(C_{2p+1})$  must contain an odd cycle whose length is obviously at most  $2p + 1$ . Thus, information regarding the shortest odd cycle of  $G_n^m$  should be useful in determining whether a homomorphism of  $C_{2p+1}$  into  $G_n^m$  can exist.

LEMMA 1. Suppose  $A, B \in V(G_n^m)$ , and suppose further that they are joined by a path of length  $2p$  ( $p \geq 0$ ) in  $G_n^m$ , then

$$|A \cap B| \geq n - (m - 2n)p.$$

*Proof.* (The author is indebted to the referee for this elegant proof.) We may assume that  $p \geq 1$ . Let  $A = X_0, X_1, \dots, X_{2p} = B$  be a path of length  $2p$  in  $G_n^m$ . This implies that  $X_i \cap X_{i+1} = \emptyset$ ,  $0 \leq i \leq 2p - 1$ . Now, for  $0 \leq i \leq p - 1$  we have

$$\begin{aligned} |X_i \cup X_{2p-i}| &\leq |(I^m - X_{i+1}) \cup (I^m - X_{2p-i-1})| \\ &= m - |X_{i+1} \cap X_{2p-i-1}| \\ &= m - 2n + |X_{i+1} \cup X_{2p-i-1}|. \\ \therefore |X_0 \cup X_{2p}| &\leq (m - 2n)(p - 1) + |X_{p-1} \cup X_{p+1}|. \end{aligned}$$



FIGURES 2, 3

Thus, we conclude

$$\begin{aligned}
 |A \cap B| &= |X_0 \cap X_{2p}| = |X_0| + |X_{2p}| - |X_0 \cup X_{2p}| \\
 &\geq 2n - (m - 2n)(p - 1) - |X_{p-1} \cup X_{p+1}| \\
 &\geq 2n - (m - 2n)(p - 1) - |I^m - X_p| \\
 &= n - (m - 2n)p.
 \end{aligned}$$

Q.E.D.

LEMMA 2. *If  $G_n^m$  contains a  $2p + 1$  cycle and if  $m \geq 2n + 1$ , then  $2p + 1 \geq m/(m - 2n)$ .*

*Proof.* Let  $A_1, A_2, \dots, A_{2p+1}$  be a  $2p + 1$  cycle in  $G_n^m$ . Then  $A_1, A_2, \dots, A_{2p+1}$  is a  $2p$  path joining  $A_1$  and  $A_{2p+1}$ .

$$\therefore |A_1 \cap A_{2p+1}| \geq n - (m - 2n)p.$$

On the other hand,  $A_1$  and  $A_{p+1}$  are also adjacent in  $G_n^m$ , so

$$\begin{aligned}
 A_1 \cap A_{2p+1} &= \emptyset. \\
 \therefore n - (m - 2n)p &\leq 0,
 \end{aligned}$$

and recalling that  $m > 2n$ , we get

$$p \geq n/(m - 2n) \quad \text{or} \quad 2p + 1 \geq m/(m - 2n).$$

Q.E.D.

It is to be noted that this lemma does not hold for even cycles. In  $G_4^9$  the vertices  $\{1, 2, 3, 4\}$ ,  $\{6, 7, 8, 9\}$ ,  $\{2, 3, 4, 5\}$ ,  $\{1, 6, 7, 8\}$ ,  $\{2, 3, 4, 9\}$ , and  $\{5, 6, 7, 8\}$  span a 6-cycle.

We next show that in the case  $m = 2n + 1$ , the lower bound on odd cycle lengths obtained in the previous lemma is actually attained.

LEMMA 3.  *$G_n^{2n+1}$  contains a  $2n + 1$  cycle.*

*Proof.* For every positive integer  $i$  we define  $A_i = \{(i - 1)n + j, 1 \leq j \leq n\}$ , where addition is mod  $2n + 1$ . Clearly, if  $i \equiv i' \pmod{2n + 1}$  then  $A_i = A_{i'}$ . The reader can easily convince himself that the converse also holds, i.e.,  $A_i = A_{i'}$  iff  $i \equiv i' \pmod{2n + 1}$ . Consequently,  $\{A_i\}$  consists of  $2n + 1$  vertices of  $G_n^{2n+1}$ . As  $A_i \cap A_{i+1} = \emptyset$  they are adjacent in  $G_n^{2n+1}$  and so  $\{A_i\}$  spans a subgraph of  $G_n^{2n+1}$  which contains a  $2n + 1$  cycle. Figure 3 is an example of such a cycle. Q.E.D.

We are now ready to derive all the chromatic numbers of  $C_{2p+1}$ .



THEOREM 6.

$$\chi_n(C_{2p+1}) = 2n + 1 + [(n-1)/p].$$

*Proof.* It is well known that for any positive integer  $n$ , there exist unique integers  $k, r$  such that

$$n - 1 = kp + r \quad k \geq 0 \quad \text{and} \quad p > r \geq 0.$$

The theorem will be proved by induction on  $k$ . We begin with  $k = 0$ . We already know that  $\chi_1(C_{2p+1}) = 3$ , and hence, by the corollary to Theorem 2 we have

$$\chi_p(C_{2p+1}) \geq 3 + 2p - 2 = 2p + 1.$$

On the other hand, since the inclusion of  $C_{2p+1}$  in  $G_p^{2p+1}$  (as given by Lemma 3 above) may be considered as a homomorphism of  $C_{2p+1}$  into  $G_p^{2p+1}$ , it follows that  $\chi_p(C_{2p+1}) \leq 2p + 1$ . Combining this with the above, we get

$$\chi_p(C_{2p+1}) = 2p + 1. \quad (8)$$

Applying the same corollary twice for each  $r$  ( $1 < r < p$ ), we get

$$\chi_p(C_{2p+1}) \geq \chi_r(C_{2p+1}) + 2p - 2r,$$

and

$$\chi_r(C_{2p+1}) \geq \chi_1(C_{2p+1}) + 2r - 2.$$

Hence, by (8),

$$\chi_r(C_{2p+1}) = 2r + 1 \quad \text{for all } 1 \leq r \leq p.$$

This completes the proof for the case  $k = 0$ .

Assume now that the theorem has been proved for  $k - 1$  and all  $0 \leq r < p$ . We will show that it also holds for  $k$  and all  $0 \leq r < p$ . Two cases are to be considered.

*Case 1.*  $r = 0$ . Here  $[(n-1)/p] = k$  and  $[(n-2)/p] = k - 1$ . Now, if it were the case that  $\chi_n(C_{2p+1}) < 2n + 1 + [(n-1)/p]$ , we would have that  $\chi_n(C_{2p+1}) \leq 2n + k = 2kp + k + 2$ . Hence, by AF1, there would exist a homomorphism

$$\theta: C_{2p+1} \rightarrow G_{kp+1}^{2pk+k+2}.$$

As we have noted before,  $\theta(C_{2p+1})$  must contain an odd cycle of length not greater than  $2p + 1$ . Lemma 2 now gives us

$$2p + 1 \geq (2kp + k + 2)/(2kp + k + 2 - 2kp - 2) = 2p + 1 + (2/k),$$

which is impossible. Hence, we must have

$$\chi_n(C_{2p+1}) \geq 2n + 1 + [(n-1)/p]. \quad (9)$$

On the other hand, it follows from Theorem 1 that  $\chi_n(C_{2p+1}) \leq \chi_1(C_{2p+1}) + \chi_{n-1}(C_{2p+1})$ . Since  $n-1 = kp$ , we have  $(n-1)-1 = (k-1)p + (p-1)$ , and we may apply the induction hypothesis to  $\chi_{n-1}(C_{2p+1})$ .

$$\begin{aligned} \therefore \chi_n(C_{2p+1}) &\leq \chi_1(C_{2p+1}) + \chi_{n-1}(C_{2p+1}) \\ &= 3 + 2(n-1) + 1 + [(n-2)/p] \\ &= 2n + 2 + k - 1 = 2n + 1 + [(n-1)/p]. \end{aligned}$$

This, together with (9) above completes the proof for the case  $r = 0$ .

*Case 2.*  $1 \leq r < p$ . Theorem 1 gives us the upper bound

$$\chi_n(C_{2p+1}) \leq \chi_{n-p}(C_{2p+1}) + \chi_p(C_{2p+1}).$$

Now,  $n-p-1 = (k-1)p + r$ , and hence, by the induction hypothesis,

$$\chi_{n-p}(C_{2p+1}) = 2(n-p) + 1 + [(n-p-1)/p] = 2(n-p) + k.$$

This and (8) now give us

$$\chi_n(C_{2p+1}) \leq 2(n-p) + k + 2p + 1 = 2n + 1 + [(n-1)/p].$$

We apply the corollary to Theorem 2 to obtain the lower bound.

$$\begin{aligned} \chi_n(C_{2p+1}) &= \chi_{pk+1+r}(C_{2p+1}) \geq \chi_{k+1}(C_{2p+1}) \\ &\quad + 2(kp+1+r) - 2(kp+1) \\ &= 2(kp+1) + 1 + [(kp+1-1)/p] + 2n - 2(kp+1) \\ &= 2n + 1 + k = 2n + 1 + [(n-1)/p]. \end{aligned}$$

So again  $\chi_n(C_{2p+1}) = 2n + 1 + [(n-1)/p]$ , and the induction is complete. Q.E.D.

The following corollary characterizes the graphs for which equality holds in Theorem 2.

**COROLLARY.** *If there exists  $n_0$  such that  $\chi_{n+1}(G) = 2 + \chi_n(G)$  for all  $n \geq n_0$ , then  $G$  is bipartite and  $\chi_{n+1}(G) = 2 + \chi_n(G)$  for all  $n$ .*

*Proof.* Clearly, if  $n \geq n_0$ , we have  $\chi_n(G) = \chi_{n_0}(G) + 2n - 2n_0$ . If

$G$  is not bipartite then it must contain an odd cycle, say  $C_{2p+1}$ . So,  $\chi_n(G) \geq \chi_n(C_{2p+1})$ . Combining these two facts we now have

$$\chi_{n_0}(G) + 2n - 2n_0 \geq 2n + 1 + [(n-1)/p]$$

or

$$\chi_{n_0}(G) - 2n_0 \geq 1 + [(n-1)/p] \quad \text{for all } n \geq n_0.$$

This, for sufficiently large  $n$ , is clearly impossible. Hence,  $G$  must be bipartite, and by Theorem 5,  $\chi_{n+1}(G) = 2 + \chi_n(G)$  for all  $n$ . Q.E.D.

Theorem 6 can be used to calculate the chromatic numbers of a variety of graphs. We demonstrate this with an example and another theorem.

EXAMPLE. Figure 4 contains 1-tuple and 2-tuple colorings of a graph  $G$ , so  $\chi_1(G) \leq 3$  and  $\chi_2(G) \leq 5$ . Since  $G$  contains a 5-cycle we must have  $\chi_n(G) \geq \chi_n(C_5) = 2n + 1 + [(n-1)/2]$ . On the other hand,

$$\chi_{2m}(G) \leq m\chi_2(G) \leq 5m = 2(2m) + 1 + [(2m-1)/2],$$

and

$$\begin{aligned} \chi_{2m+1}(G) &\leq m\chi_2(G) + \chi_1(G) \leq 5m + 3 \\ &= 2(2m+1) + 1 + [(2m+1-1)/2]. \end{aligned}$$

And combining the above inequalities we have

$$\chi_n(G) \leq 2n + 1 + [(n-1)/2].$$

Since the reverse inequality was derived before, we have

$$\chi_n(G) = 2n + 1 + [(n-1)/2].$$

The same method can be applied to derive the chromatic numbers of  $G_m^{2m+1}$ .

THEOREM 7.  $\chi_n(G_m^{2m+1}) = 2n + 1 + [(n-1)/m]$ .

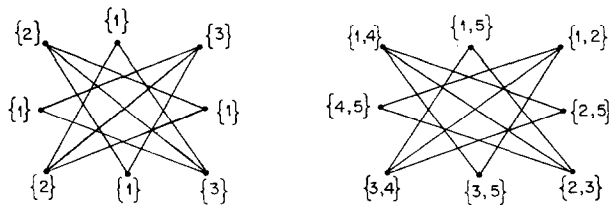


FIGURE 4

*Proof.* We already know from Lemma 3 of Theorem 6 that  $G_m^{2m+1}$  contains a  $2m + 1$  cycle. Consequently we have the lower bound

$$\chi_n(G_m^{2m+1}) \geq \chi_n(C_{2m+1}) = 2n + 1 + [(n - 1)/m]. \quad (10)$$

The reverse inequality will again be proved by induction on  $k$  where  $n - 1 = km + r$ ,  $k \geq 0$ ,  $m > r \geq 0$ . Suppose  $k = 0$ . By Theorem 2,  $\chi_n(G_m^{2m+1}) \leq \chi_m(G_m^{2m+1}) - 2(m - n)$ ,  $1 \leq n \leq m$ . Since the identity map  $G_m^{2m+1} \rightarrow G_m^{2m+1}$  is a homomorphism,  $\chi_m(G_m^{2m+1}) \leq 2m + 1$ . This completes the proof for the case  $k = 0$  and  $1 \leq n \leq m$ . Assume now that the theorem holds for  $k - 1$ . As  $n - m - 1 = (k - 1)m + r$ , we may apply the induction hypothesis to  $n - m$ , so that we have

$$\begin{aligned} \chi_n(G_m^{2m+1}) &\leq \chi_{n-m}(G_m^{2m+1}) + \chi_m(G_m^{2m+1}) \\ &= 2(n - m) + 1 + [(n - m - 1)/m] + 2m + 1 \\ &= 2n + 1 + [(n - 1)/m]. \end{aligned} \quad \text{Q.E.D.}$$

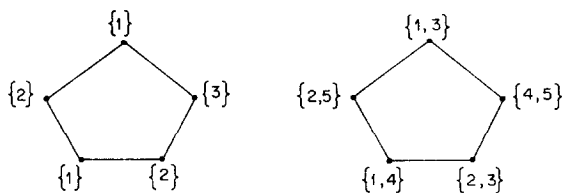
#### 4. EFFICIENT COLORINGS

Figures 5 and 6 below show examples of minimal 1-tuple and 2-tuple colorings of  $C_5$ . The color classes in Fig. 5 consist of two maximal independent subsets and a singleton. The color classes in Fig. 6 consist of five maximal independent subsets. Thus we might say that the 2-tuple coloring makes "better" use of the independent subsets of  $C_5$ . This notion is made precise by the following definition.

**DEFINITION.** An  $n$ -tuple coloring of  $G$  with  $m$  colors is said to be *efficient* if

$$m/n \leq (\chi_k(G))/k \quad \text{for all } k \geq 1.$$

It is clear that any efficient coloring is necessarily minimal. On the other hand, it is not known whether every graph has an efficient coloring. Theorem 8 gives a sufficient condition for the existence of such efficient



FIGURES 5, 6

colorings. We first recall that an independent set is said to be *maximum* if its cardinality is not less than that of any other independent subset of the same graph.

**THEOREM 8.** *If  $\{C_\lambda\}$  is an *n*-tuple coloring of *G* with *m* colors, and if each  $C_\lambda$  is a maximum independent subset of  $V(G)$ , then  $\{C_\lambda\}$  is an efficient coloring of *G*. Moreover, all efficient colorings of *G* must then consist of maximum independent sets.*

*Proof.* Let  $\beta_0 = |C_\lambda|$  for all  $\lambda$ ,  $p = |V(G)|$  and suppose  $D_\nu$  is a minimal *k*-tuple coloring of *G*. It follows from AF2 that

$$\sum_{\lambda=1}^m |C_\lambda| = pn \quad \text{and} \quad \sum_{\nu=1}^{x_k(G)} |D_\nu| = pk.$$

Since  $\beta_0 = |C_\lambda| \geq |D_\nu|$  for all  $\lambda, \nu$ , we have

$$\beta_0 m = pn \quad \text{and} \quad \beta_0 \chi_k(G) \geq pk \quad (11)$$

or

$$m/n = p/\beta_0 \leq (\chi_k(G))/k.$$

Hence  $\{C_\lambda\}$  is an efficient coloring. Moreover,  $m/n = (\chi_k(G))/k$  iff equality holds in (11); in other words iff each  $D_\nu$  is maximum. Q.E.D.

Before we proceed to derive the efficient colorings of the  $G_n^m$ , we must digress and examine the independent subsets of  $V(G_n^m)$  closely. We know that two vertices of  $G_n^m$  are not adjacent iff they are not disjoint. Hence, an independent set of vertices of  $G_n^m$  consists of a family  $\{a_\lambda\}$  of pairwise nondisjoint subsets of  $I^m$  such that  $|a_\lambda| = n$  for all  $\lambda$ . Such families have already been investigated by other authors (see [1, 4]) and are said to have the *intersection property*. We follow the notation of [4] where  $S(1, n, m)$  denotes, in effect, the family of all independent sets of vertices of  $G_n^m$ , and write,

**THEOREM OF ERDÖS-CHAO KO-RADO.** *If  $1 \leq n \leq \frac{1}{2}m$  and*

$$\{a_\lambda\} \in S(1, n, m), \quad \text{then} \quad |\{a_\lambda\}| \leq \binom{m-1}{n-1}.$$

Now [4, Remark 2, p. 334] is not quite valid, as equality can hold even if  $m > 2l - k$ . This can happen if  $\{a_\lambda\}$  does not happen to contain all the available integers, i.e., if  $\bigcup_\lambda a_\lambda \neq \{0, 1, 2, \dots, m-1\}$ . The crucial point in the proof of [4, Theorem 2] occurs on p. 332, the sentence "If  $n_0 = 0, \dots$ " For by applying the induction hypothesis we may carry over equality to

the case  $m > 2l - k$ . This point, minor as it may be, enables us to strengthen the above theorem.

**THEOREM OF ERDÖS-CHAO KO-RADO (Strong Version).** *If  $1 \leq n \leq \frac{1}{2}m$  and  $\{a_\lambda\} \in S(1, n, m)$ , then  $|\{a_\lambda\}| \leq \binom{m-1}{n-1}$  and if equality holds and  $m > 2n$ , then  $|\bigcap_\lambda a_\lambda| = 1$ .*

*Proof.* We will merely indicate the required additions to Katona's proof of [4, Remark 3, p. 334]. If  $|\{a_\lambda\}| = \binom{m-1}{n-1}$ , then equality must hold in (10). Hence, by the corrected version of Remark 2 given above, either  $m = 2(m-l) - (m-2l+1)$ , or  $m-n=n$ , or  $\bigcup_\nu b_\nu \neq \{1, 2, \dots, m\}$ . The first alternative is manifestly self-contradictory and the second contradicts  $m > 2n$ . Consequently, the third alternative must hold.

$$\therefore \bigcap_\lambda a_\lambda \neq \emptyset.$$

However, if  $|\bigcap_\lambda a_\lambda| = k$ , then  $|\{a_\lambda\}| \leq \binom{m-k}{n-k}$ , and hence  $k = 1$ .

Q.E.D.

It now follows easily from the strengthened version of the Erdős-Chao Ko-Rado theorem that the maximum independent sets of vertices of  $G_n^m$  are the sets  $C_i$  where  $C_i = \{A \in V(G_n^m) \mid i \in A\}$ . The derivation of the efficient colorings of  $G_n^m$  is now quite an easy matter.

**THEOREM 9.** *If  $m > 2n$ , then  $G_n^m$  has an efficient  $k$ -tuple coloring iff  $n$  divides  $k$ .*

*Proof.* Let  $C_i$  be as defined above. Clearly  $\{C_i\}$  for  $1 \leq i \leq m$  is an  $n$ -tuple coloring of  $G_n^m$ . By Theorem 8 this coloring is efficient. Moreover, writing  $\chi_i$  for  $\chi_i(G_n^m)$ ,

$$\begin{aligned} \chi_n/n = m/n \leq \chi_{sn}/sn \leq s\chi_n/sn = \chi_n/n \quad \text{for all } s \geq 1. \\ \therefore \chi_{sn}/sn = \chi_n/n = m/n. \end{aligned} \tag{12}$$

In other words, if  $n$  divides  $k$ , then  $G_n^m$  has an efficient  $k$ -tuple coloring. Conversely, suppose  $G_n^m$  has a  $k$ -tuple efficient coloring  $\{D_\lambda\}$ . By Theorem 8, each  $D_\lambda$  is a maximum independent set. By the strong Erdős-Chao Ko-Rado theorem, for each  $D_\lambda$  there is an integer  $j_\lambda$  such that  $j_\lambda \in A$  for all  $A \in D_\lambda$ . Let  $\nu_j$  be the number of  $D_\lambda$  for which  $j_\lambda = j$ . Now, a vertex  $A$  of  $G_n^m$  is contained in  $D_\lambda$  iff  $j_\lambda \in A$ . Since each vertex is contained in  $k$  of the  $D_\lambda$ , it follows that

$$\sum_{j \in A} \nu_j = k.$$

For any two integers  $h$  and  $i$ , where  $1 \leq h \neq i \leq n$ , choose vertices  $A$  and  $B$  so that they differ only by  $h$  and  $i$ ; in other words, choose  $A$  and  $B$  so that  $h \in A$ ,  $i \in B$  and  $A = B - \{i\} + \{h\}$ . Then

$$0 = k - k = \left( \sum_{j \in A} v_j \right) - \left( \sum_{s \in B} v_s \right) = v_h - v_i.$$

So,

$$v_h = v_i \quad \text{for all } 1 \leq h, i \leq n.$$

Hence, if  $v$  is the common value of all the  $v_j$ , then each vertex of  $G_n^m$  is contained in  $nv$  of the  $D_\lambda$ .

$$\therefore k = nv \quad \text{and} \quad n \text{ divides } k.$$

Q.E.D.

COROLLARY.  $\chi_{kn}(G_n^m) = km$ .

*Proof.* See (12).

Q.E.D.

Efficient colorings can be used to derive the chromatic numbers of nonsymmetric graphs as well. Consider, for example, the graph  $G$  of Fig. 7. Clearly  $\chi_1(G) = 4$ . The reader can easily convince himself, applying Theorem 8, that the 2-tuple coloring given in Fig. 7 is in fact efficient. Hence, by Theorem 9,  $\chi_{2n}(G) = 7n$ . Now, by Theorem 1,  $\chi_{2n+1}(G) \leq \chi_1(G) + \chi_{2n}(G) = 4 + 7n$ . Since the 2-tuple coloring is efficient we must also have

$$\begin{aligned} (\chi_{2n+1}(G))/(2n+1) &\geq 7/2. \\ \therefore \chi_{2n+1}(G) &\geq 7n + (7/2). \end{aligned}$$

And since  $\chi_{2n+1}(G)$  is an integer we must have  $\chi_{2n+1}(G) = 7n + 4$ .

Theorems 8 and 9 can be used to calculate some more of the numbers  $\chi_k(G_n^m)$ . However, before proceeding further with these calculations, the reader should be made aware of the significance of these numbers. This

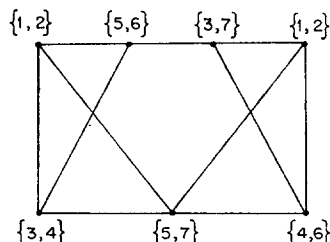


FIGURE 7

can best be done with the aid of a concrete example. It follows from Theorem 7 that  $\chi_6(G_4^9) = 14$ . Hence, there exists a homomorphism  $G_4^9 \rightarrow G_9^{14}$ . In view of AF1 it now follows that if any graph can be given a 4-tuple coloring with nine colors, then at most 14 colors are required for a 6-tuple coloring. Thus the numbers  $\chi_k(G_n^m)$  contain valuable information relating different chromatic numbers. The proof of Theorem 10 below serves as an example for the utilization of this information. First, however, a lemma is needed.

LEMMA. *If  $\chi_k(G_n^m) = \chi_k(G_t^s)$  for all  $k$ , then  $m = s$  and  $n = t$ .*

*Proof.* Since  $G_n^m$  has an  $n$ -tuple efficient coloring and  $G_t^s$  has a  $t$ -tuple efficient coloring, it follows from Theorem 9 that  $n$  divides  $t$ , and  $t$  divides  $n$ .

$$\therefore n = t.$$

Moreover,

$$m = \chi_n(G_n^m) = \chi_n(G_t^s) = \chi_t(G_t^s) = s.$$

Q.E.D.

THEOREM 10.  $\chi_k(G_n^{2n+2}) = 2 + 2k$  if  $(n-1)/2 < k \leq n$ .

*Proof.* If  $k = n$ , then this theorem reduces to a special case of the corollary to Theorem 9. Hence, we may assume that  $(n-1)/2 < k < n$ , or,

$$[(n-1)/k] = 1.$$

By Theorem 7,  $\chi_n(G_k^{2k+1}) = 2n + 1 + [(n-1)/k] = 2n + 2$ , and so there exists a homomorphism

$$G_k^{2k+1} \rightarrow G_n^{2n+2}. \quad (13)$$

Suppose now that  $\chi_k(G_n^{2n+2}) \leq 2k + 1$ , then there exists a homomorphism

$$G_n^{2n+2} \rightarrow G_k^{2k+1}. \quad (14)$$

It now follows from Theorem 3 and the lemma (applied to (13) and (14)) that  $2n + 2 = 2k + 1$  and  $n = k$ , which is absurd. Hence, we must have  $\chi_k(G_n^{2n+2}) \geq 2k + 2$ . The upper bound is given to us by the corollary to Theorem 2:

$$2n + 2 = \chi_n(G_n^{2n+2}) \geq \chi_k(G_n^{2n+2}) + 2n - 2k,$$

or

$$\chi_k(G_n^{2n+2}) \leq 2k + 2. \quad \text{Q.E.D.}$$



A function  $f: I^+ \rightarrow I^+$  is said to be *subadditive* if  $f(a + b) \leq f(a) + f(b)$ . Theorem 1 states that  $\chi_n(G)$  is a subadditive function of its argument  $n$ . The following theorem on subadditive functions again demonstrates the significance of efficient colorings.

**THEOREM 11.** *If  $f: I^+ \rightarrow I^+$  is a subadditive function, and if there exists  $n$  such that*

$$f(x)/x \geq f(n)/n \quad \text{for all } x \geq 1,$$

*then there exists an integer  $x_0$  such that*

$$f(x + n) = f(x) + f(n) \quad \text{for all } x \geq x_0.$$

*Proof.* Since  $f$  is subadditive we know that  $f(kn) \leq kf(n)$ . But we are also given that

$$\begin{aligned} f(kn)/kn &\geq f(n)/n. \\ \therefore f(kn) &= kf(n). \end{aligned}$$

Now, for arbitrary  $k$  and  $0 \leq r < n$ , define

$$\Delta(k, r) = f(kn + r) - f(kn).$$

Then

$$\begin{aligned} \Delta(k, r) - \Delta(k + 1, r) &= f(kn + r) - f(kn) - f((k + 1)n + r) \\ &\quad + f((k + 1)n) \\ &= f(n) + f(kn + r) - f(kn + r + n) \geq 0. \end{aligned} \tag{15}$$

The last part of the above inequality is due to the subadditivity of  $f$ . On the other hand,

$$\begin{aligned} f(kn + r) + f(n - r) &\geq f((k + 1)n) = (k + 1)f(n) = kf(n) + f(n) \\ &= f(kn) + f(n) \\ \therefore f(kn + r) - f(kn) &\geq f(n) - f(n - r), \end{aligned}$$

or,

$$\Delta(k, r) \geq f(n) - f(n - r) \quad \text{for all } k.$$

Hence,

$$\Delta(k, r) \geq \Delta(k + 1, r) \geq f(n) - f(n - r) \quad \text{for all } k,$$

and as  $\Delta(k, r)$  can assume only integral values, it follows that there exists

an integer  $m_r$  such that  $\Delta(k, r) = \Delta(k + 1, r)$  whenever  $k \geq m_r$ . Let  $M = \text{Max}\{m_r\}$ , then,

$$\text{for } k \geq M \quad \Delta(k, r) = \Delta(k + 1, r) \quad \text{for all } 0 \leq r < n.$$

Set  $x_0 = Mn$ . If  $x \geq x_0$  then there exist integers  $k$  and  $r$ , satisfying  $k \geq M$  and  $0 \leq r < n$ , such that  $x = kn + r$ . Consequently,  $\Delta(k, r) = \Delta(k + 1, r)$ , and equality must hold in (15).

$$\therefore f(n) + f(kn + r) - f(kn + r + n) = 0.$$

In other words,

$$f(x) + f(n) = f(x + n).$$

Q.E.D.

**COROLLARY.** For a fixed  $G_n^m$  and sufficiently large  $k$ ,

$$\chi_{k+n}(G_n^m) = m + \chi_k(G_n^m).$$

*Proof.* Merely apply Theorems 9 and 11 to  $G_n^m$ .

Q.E.D.

## 5. CONCLUSION

A pattern for  $\chi_k(G_n^m)$  emerges from the theorems of this paper and several other calculations not displayed here. It would seem that in general  $\chi_1(G_n^m) = m - 2n + 2$ , and consequently, by Theorem 2,

$$\chi_k(G_n^m) = 2 + \chi_{k-1}(G_n^m) \quad \text{for } 1 < k \leq n.$$

For  $k = n + 1$  it seems again that  $\chi_{n+1}(G_n^m) = \chi_n(G_n^m) + \chi_1(G_n^m)$ . Again, since  $\chi_{2n}(G_n^m) = 2m$ , this would entail

$$\chi_k(G_n^m) = 2 + \chi_{k-1}(G_n^m) \quad \text{for } n + 1 < k \leq 2n.$$

We generalize this to the following.

*Conjecture.* If  $k - 1 = qn + r$  where  $q \geq 0$  and  $0 \leq r < n$ , then

$$\chi_k(G_n^m) = (q + 1)m - 2(n - r - 1).$$

*Added in Proof.* The author wishes to thank the referee for the current version of Theorem 11, as well as for many other helpful suggestions. In addition, it should be pointed out that it has recently been proved in [5] and [6] that every graph does indeed possess efficient colorings. In fact, a slightly different version of Theorem 11 is known to Scott [5], whose as yet unpublished results predate those presented here.

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